

Quantum Computing

- Week 4 (May 5-6, 2026)
- Today:
 - Postulates of Quantum Computing

Postulates of Quantum Computing

“Don’t be surprised if the motivation for the postulates is not always clear; even to experts the basic postulates of quantum mechanics appear surprising...” from [NC00]

Postulates of Quantum Computing

- First postulate: State space

Postulate 1: Associated to any isolated physical system is a complex vector space with inner product (that is, a Hilbert space) known as the *state space* of the system. The system is completely described by its *state vector*, which is a unit vector in the system's state space. from [NC00]

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from [NC00]

- Keywords:
 - **Isolated system**
 - **Hilbert space:** Complex inner product linear space (e.g., \mathbb{C}^{2^n})
 - The state of a system is completely described by a **state vector**
 - A state vector is a **unit vector of** the Hilbert space

Example:
 $|\phi\rangle = \alpha|0\rangle + \beta|1\rangle$

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- Keywords:
 - **Isolated system: (Informally,) Not entangled with other systems...**
 - **Hilbert space:** Complex inner product linear space (e.g., \mathbb{C}^{2^n})
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Postulates of Quantum Computing

- Second postulate: Evolution

Postulate 2: The evolution of a *closed* quantum system is described by a *unitary transformation*. That is, the state $|\psi\rangle$ of the system at time t_1 is related to the state $|\psi'\rangle$ of the system at time t_2 by a unitary operator U which depends only on the times t_1 and t_2 ,

$$|\psi'\rangle = U|\psi\rangle .$$

from [NC00]

- Keywords:
 - **Closed system**
 - **Unitary transformation**

Postulates of Quantum Computing

- Second postulate (using Schrodinger's equation): Evolution

Postulate 2': The time evolution of the state of a closed quantum system is described by the *Schrödinger equation*,

$$i\hbar \frac{d|\psi\rangle}{dt} = H|\psi\rangle.$$

In this equation, \hbar is a physical constant known as *Planck's constant* whose value must be experimentally determined. The exact value is not important to us. In practice, it is common to absorb the factor \hbar into H , effectively setting $\hbar = 1$. H is a fixed Hermitian operator known as the *Hamiltonian* of the closed system.

from [NC00]

Postulates of Quantum Computing

- Fourth postulate: Composite system:

Postulate 4: The state space of a composite physical system is the tensor product of the state spaces of the component physical systems. Moreover, if we have systems numbered 1 through n , and system number i is prepared in the state $|\psi_i\rangle$, then the joint state of the total system is $|\psi_1\rangle \otimes |\psi_2\rangle \otimes \cdots \otimes |\psi_n\rangle$. from [NC00]

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- Keywords:
 - (Suppose each component system is **isolated**)
 - **Tensor product**
 - State space of the whole system: **Tensor product of the Hilbert spaces** of each component system
 - State vector of the whole system: **Tensor product of the state vector** of each component system

Postulates of Quantum Computing

- Third postulate: Quantum measurement
- Do it on the board
- Measurement **in the computational basis**
- **Partial measurement**

Postulates of Quantum Computing

- Third postulate: Quantum measurement
- Measurement in the computational basis
- Partial measurement
- **Collapse: The state after measurement**

$$|\phi\rangle \longrightarrow \frac{M_m|\phi\rangle}{\sqrt{\langle\phi|M_m^\dagger M_m|\phi\rangle}}$$

Quantum Measurement

- Let $\{M_m\}_m$ be a set of matrices describing some quantum measurement
- Let $|\phi\rangle$ be a quantum state, perform the same measurement $\{M_m\}_m$ on $|\phi\rangle$ **twice**

$$|\phi\rangle \longrightarrow \frac{M_m|\phi\rangle}{\sqrt{\langle\phi|M_m^\dagger M_m|\phi\rangle}} \longrightarrow ?$$

(with probability
 $p(m) = \langle\phi|M_m^\dagger M_m|\phi\rangle$)

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We need more restrictions to achieve
“Stability”

Projective Measurement

- Projective measurements: A special class of measurements

Projective Measurement

- Some Linear Algebra – **Projector and Eigenspace:**
 - A matrix P is a projector (or projection operator) if $P^2 = P$
 - For any vector x , $P^n x = P^{n-1} x = \dots = Px$
 - **Eigenvalues and Eigenvectors:** $Ax = \lambda x$
 - A matrix (linear operator) may have **multiple eigenvalues**
 - Each eigenvalue may have **multiple linearly independent eigenvectors**
 - Let $\{x_1, \dots, x_m\}$ denote a *maximal set* of linearly independent eigenvectors of λ (i.e., $Ax_i = \lambda x_i$)
 - We say $\{x_1, \dots, x_m\}$ **span the eigenspace** of A with eigenvalue λ
 - Given such $\{x_1, \dots, x_m\}$, the **Gram-Schmidt process** gives us an **orthogonal basis of the eigenspace**
 - Given an orthogonal basis, we can easily compute the **projector onto this eigenspace**

Projective Measurement

- Some Linear Algebra – **Spectral decomposition (simplified)**:

Any Hermitian operator M (i.e., $M = M^\dagger$) can be written as:

$$M = \sum_{\lambda} \lambda P_{\lambda}$$

- λ represents an eigenvalue of M
- P_{λ} represents the projector onto the λ eigenspace
- P_{λ} itself is also Hermitian, i.e., $P_{\lambda} = P_{\lambda}^\dagger$
- Examples (show on the board)...

Projective Measurement

- Projective measurements: A special class of measurements
- (Do it on the board)
- Keywords
 - **Observable** $M = \sum_m mP_m$ is a *Hermitian matrix*
 - m represents an eigenvalue of M , and it is also used to **label a measurement outcome**
 - P_m represents the projector onto the m eigenspace
 - P_m is also Hermitian
 - Measurement outcome **correspond to the eigenvalues**
 - e.g., $p(m) = \langle \phi | P_m | \phi \rangle$
 - The state after measurement: $|\phi\rangle \rightarrow \frac{P_m|\phi\rangle}{\sqrt{p(m)}}$

Projective Measurement

- Projective measurements: A special class of measurements
- Relation to Postulate 3:
 - $M = \sum_m mP_m$, but $\sum_m P_m = I$, so the completeness condition holds
 - **Note: An eigenvalue of an observable just represents a possible outcome (i.e., a label), but not the probability or physical meaning by itself**
- Examples: $|0\rangle\langle 0| = 1 \cdot |0\rangle\langle 0| + 0 \cdot |1\rangle\langle 1|$ and $|1\rangle\langle 1| = 0 \cdot |0\rangle\langle 0| + 1 \cdot |1\rangle\langle 1|$
 - Both can be used to represent measurement in the computational basis

Projective Measurement

- Let $M = \sum_m mP_m$ be an observable
- Let $|\phi\rangle$ be a quantum state, perform the same projective measurement M on $|\phi\rangle$ **twice**

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$$|\phi\rangle \longrightarrow \frac{P_m|\phi\rangle}{\sqrt{\langle\phi|P_m|\phi\rangle}} \longrightarrow \frac{P_m|\phi\rangle}{\sqrt{\langle\phi|P_m|\phi\rangle}}$$

(with probability $p(m) = \langle\phi|P_m|\phi\rangle$)

(with probability 1)

Non-Projective Measurement

- Let $\{M_m\}_m$ be a set of matrices describing some quantum measurement
- General measurement (Postulate 3): $M_m^\dagger M_m$ is not necessarily a projector.
 - **Non-projective measurement:** $M_m^\dagger M_m$ is not a projector
 - Cannot guarantee that the same result will be reproduced if the same measurement is repeated
- Used in various quantum information-processing protocols
 - But will not be covered in this course

The Deutsch-Jozsa Problem

- Constant-vs-balanced problem
- Let $f: \{0,1\}^n \rightarrow \{0,1\}$ be a bit function such that it is in either two cases:
 - f is a *constant* function: $\forall x \in \{0,1\}^n, f(x)$ is always a constant (0 or 1)
 - f is a *balanced* function: $\sum_{x \in \{0,1\}^n} f(x) = 2^{n-1}$ (i.e., outputs 0 for half the inputs, and 1 for the other half)
- To decide whether f is constant or balanced, **how many times** must we evaluate f ?

Classical Computer

Worst-case: 2^n

Probabilistic algorithm:

$l \ll 2^n$ times,

with a failure rate of $\frac{1}{2^l}$

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Classical Computer

Worst-case: 2^n

Probabilistic algorithm:

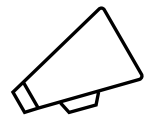
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Can we do better?

References

- [NC00]: Sections 1.4.4, 2.2
- [KLM07]: Chapter 3, Sections 6.4.



No lectures next week (May 12-13)